

# Random Continuous Functions

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## Goals

**Extend the notion of Martin-Löf random reals to**

1. Random closed sets
2. Random continuous functions
3. Random structures.

## ML Random Reals

A real  $x \in 2^{\mathbb{N}}$  is **Martin-Löf random** if  
for any effective sequence  $S_1, S_2, \dots$  of c.e. open sets with  $\mu(S_n) \leq 2^{-n}$ ,

$$x \notin \bigcap_n S_n.$$

## Closed Sets and Trees

For a finite string  $\sigma$ , let  $I(\sigma)$  denote  $\{x \in 2^{\mathbb{N}} : \sigma \prec x\}$ .

We shall call  $I(\sigma)$ , the **interval** determined by  $\sigma$ .

A nonempty closed set  $P$  may be identified with a tree  $T_P \subseteq \{0, 1\}^*$  where

$$T_P = \{\sigma : P \cap I(\sigma) \neq \emptyset\}.$$

Note that  $T_P$  has no dead ends, i.e., if  $\sigma \in T_P$ ,

then either  $\sigma \frown 0 \in T_P$  or  $\sigma \frown 1 \in T_P$ .

## Closed Sets and Trees II

For an arbitrary tree  $T \subseteq \{0, 1\}^*$ , let  $[T]$  denote the set of infinite paths through  $T$ , that is,

$$x \in [T] \iff (\forall n)x \upharpoonright n \in T.$$

$P \subseteq 2^{\mathbb{N}}$  is a closed set if and only if  $P = [T]$  for some tree  $T$ .

$P$  is a  $\Pi_1^0$  class, or an effectively closed set, if  $P = [T]$  for some computable tree  $T$ .

The complement of a  $\Pi_1^0$  class is called a c.e. open set.

There is a natural effective enumeration  $P_0, P_1, \dots$  of the  $\Pi_1^0$  classes and thus an enumeration of the c.e. open sets.

## Random Closed Sets

Due to Broadhead, Cenzer, and Dashti

Given a closed set  $Q \subseteq 2^{\mathbb{N}}$ , let  $T = T_Q$  be the tree without dead ends such that  $Q = [T]$ .

Let  $\sigma_0, \sigma_1, \dots$  enumerate the elements of  $T$  in order, first by length and then lexicographically.

We then define the **code**  $x = x_Q = x_T$  by recursion such that for each  $n$ ,

- (i)  $x(n) = 2$  if both  $\sigma_n \frown 0$  and  $\sigma_n \frown 1$  are in  $T$ ,
- (ii)  $x(n) = 1$  if  $\sigma_n \frown 0 \notin T$  and  $\sigma_n \frown 1 \in T$ , and
- (iii)  $x(n) = 0$  if  $\sigma_n \frown 0 \in T$  and  $\sigma_n \frown 1 \notin T$ .

We then define a **measure**  $\mu^*$  on  $\mathcal{C}$  by setting

$$\mu^*(\mathcal{X}) = \mu(\{x_Q : Q \in \mathcal{X}\}) \quad (1)$$

for any  $\mathcal{X} \subseteq \mathcal{C}$  where  $\mu$  is the standard measure on  $\{0, 1, 2\}^{\mathbb{N}}$ .

## Random Closed Sets II

Thus each of the following scenarios has equal probability  $\frac{1}{3}$ :

$\sigma \frown 0 \in T$  and  $\sigma \frown 1 \in T$ ,

$\sigma \frown 0 \in T$  and  $\sigma \frown 1 \notin T$ , and

$\sigma \frown 0 \notin T$  and  $\sigma \frown 1 \in T$ .

**Definition.** A closed set  $Q \subseteq 2^{\mathbb{N}}$  is **(Martin-Löf) random** if  $x_Q$  is (Martin-Löf) random.

## Random Closed Sets III

**Theorem 1** (BCD 2006) *There is a  $\Pi_2^0$  random closed set.*

(There are  $\Delta_2^0$  random reals)

**Theorem 2** (BCD 2006) *If  $Q$  is a random closed set, then  $Q$  has no computable members.*

**Theorem 3** (BCD 2006) *If  $Q$  is a random closed set, then  $Q$  has no isolated members.*

**Theorem 4** (BCD 2006) *If  $Q$  is a random closed set, then  $\mu^*(Q) = 0$ .*

## Random Closed Sets IV

**Theorem 5** (BCD 2006) *If  $Q$  is a  $\Pi_1^0$  class with measure 0, then no subset of  $Q$  is a random closed set.*

**Corollary 6** *No  $\Pi_1^0$  class is a random closed set.*

## Random Continuous Functions

A continuous function on  $2^{\mathbb{N}}$  is a function with a closed graph.

Thus we might say that a function  $F$  is random if the graph

$Gr(F) = \{x \oplus y : y = F(x)\}$  is a random closed set.

**Problem:** If  $[T]$  is the graph of a function and  $\sigma \in T$  has even length, then we must have  $\sigma \frown 0 \in T$  and  $\sigma \frown 1 \in T$ .

Thus the family of closed sets which are the graphs of functions has measure 0 in the space of closed sets and hence a random closed set will not be the graph of a function.

## Representing Functions

A continuous function  $F : 2^{\mathbb{N}} \rightarrow 2^{\mathbb{N}}$  may be represented by a function  $f : \{0, 1\}^* \rightarrow \{0, 1\}^*$  such that the following hold for all  $\sigma \in \{0, 1\}^*$ .

- (1)  $|f(\sigma)| \leq |\sigma|$ .
- (2)  $\sigma_1 \prec \sigma_2$  implies  $f(\sigma_1) \preceq f(\sigma_2)$ .
- (3) For every  $n$ , there exists  $m$  such that for all  $\sigma \in \{0, 1\}^m$ ,  $|f(\sigma)| \geq n$ .
- (4) For all  $x \in 2^{\mathbb{N}}$ ,  $F(x) = \bigcup_n f(x \upharpoonright n)$ .

We define the space  $\mathcal{F}$  of representing functions  $f : \{0, 1\}^* \rightarrow \{0, 1\}^*$  to be those which satisfy clauses (1) and (2) above.

## Properties of Representing Functions

- 1) If  $f$  is a representing function of  $F$ ,  $f(\emptyset) = \emptyset$  by (1).
- 2) If  $f((0)) = (i)$  where  $i \in \{0, 1\}$ , then for all  $x \in I((0))$ ,  $F(x)(0) = i$ .
- 3) If  $f((0)) = \emptyset$ , then we shall take this to mean that there exist  $x_0$  and  $x_1$  in  $I((0))$  such that  $F(x_i)(0) = i$  for  $i = 0, 1$ .
- 4)  $F(\sigma) \preceq \tau$  where  $\tau$  is the longest string  $\tau$  with  $|\tau| \leq |\sigma|$  such that  $\tau \prec F(x)$  whenever  $\sigma \prec x$ .

## Coding Representing Functions

There is a one-to-one correspondence between  $\mathcal{F}$  and  $\{0, 1, 2\}^{\mathbb{N}}$  defined as follows.

Enumerate  $\{0, 1\}^*$  in order, first by length and then lexicographically, as  $\sigma_0, \sigma_1, \dots$ . Thus  $\sigma_0 = \emptyset$ ,  $\sigma_1 = (0)$ ,  $\sigma_2 = (1)$ ,  $\sigma_3 = (00)$ ,  $\dots$

Then  $r \in \{0, 1, 2\}^{\mathbb{N}}$  corresponds to the function  $f_r : \{0, 1\}^* \rightarrow \{0, 1\}^*$  defined by declaring that  $f_r(\emptyset) = \emptyset$  and that, for any  $\sigma_n$  with  $|\sigma_n| \geq 1$ ,

$$f_r(\sigma_n) = \begin{cases} f_r(\sigma_k), & \text{if } r(n) = 2; \\ f_r(\sigma_k) \frown i, & \text{if } r(n) = i < 2. \end{cases}$$

where  $k$  is such that  $\sigma_n = \sigma_k \frown j$  for some  $j$ .

Then every continuous function  $F$  has a representative  $f$  as described above, and, in fact, it has infinitely many representatives.

## The measure $\mu^{**}$

The measure  $\mu^{**}$  on  $\mathcal{F}$  is then induced by the standard probability measure on  $\{0, 1, 2\}^{\mathbb{N}}$ .

**Definition.** A **effectively random continuous function** on  $2^{\mathbb{N}}$  is one which has a representation in  $\mathcal{F}$  which is Martin-Löf random.

**Remark:** Fouche (2000) has a different approach to randomness for continuous functions connected with Brownian motion, first presented by Asarin and Prokovsky (1986).

## Properties of Random Functions

Our first result will be to show that every random function represents a continuous function. To prove such a result, we need to prove the following lemma.

**Lemma 7** *Let  $\Sigma$  be a finite set and let  $Q \subseteq \Sigma^{\mathbb{N}}$  be a  $\Pi_1^0$  class of measure 0. Then no element of  $Q$  is Martin-Löf random.*

**Theorem 8** *The set of functions in  $\mathcal{F}$  which represent a total continuous function has measure one. Hence every random function represents a continuous function.*

**Theorem 9** *There exists a random continuous function which is  $\Delta_2^0$  computable.*

## Properties of Random Functions II

**Proposition 10** *If  $F$  is a random continuous function, then, for every  $\sigma \in \{0, 1\}^*$ , the function  $F_\sigma$  is random continuous, where*

$$F_\sigma(x) = F(\sigma \frown x).$$

**Theorem 11** *If  $F$  is a random continuous function, then, for any computable real  $x$ ,  $F(x)$  is a random real.*

**Remark:** Fouche has shown that for his definition of random continuous function, any random continuous function  $F$ ,  $F(x)$  is not computable for any computable input  $x$ .

**Theorem 12** *If  $F$  is a random continuous function, then the set of zeros of  $F$  is a random closed set.*

## Properties of Random Functions III

**Theorem 13** *If  $F$  is a random continuous function, then the image  $F[2^{\mathbb{N}}]$  has no isolated elements.*

**Theorem 14** *For any  $\sigma \in \{0, 1\}^*$ , the probability that the image of a continuous function  $F$  meets  $I(\sigma)$  is always  $> \frac{3}{4}$ .*

**Corollary 15** *The image of a random continuous function need not be a random closed set.*

## Random Structures

**Idea** Use random subsets of a universal structure.

**(Martin-Löf) Random Graphs.** Start with  $K_{\mathbb{N}}$  which is the complete graph on  $\mathbb{N}$ .

Let  $e_0, e_1, \dots$  be an effective list of all edges of  $K_{\mathbb{N}}$ .

Then given a real  $r = r_0 r_1 \dots \in \{0, 1\}^{\mathbb{N}}$ , we let  $K_r$  be the graph  $(\mathbb{N}, E_r)$  where

$$e_i \in E_r \iff r_i = 1.$$

**Definition** We say that  $K_r$  is **(Martin-Löf) random** if and only if  $r$  is (Martin-Löf) random.

## (Martin-Löf) Random Graphs

We say that  $K_r$  has **property**  $P_{n,s}$  if for any subset  $V = \{v_1, \dots, v_n\}$  of vertices of  $K_r$  of size  $n$  and for any subset  $S \subseteq V$  of size  $s$ , there is a vertex  $w \notin V$  such that for all  $i = 1, \dots, n$ ,

$$\{v_i, w\} \in E_r \iff v_i \in S.$$

Thus  $w$  witnesses that there is a vertex which is a neighbor of each vertex in  $S$  and is not a neighbor of each vertex in  $V - S$ .

## Random Graphs

It is easy to see by a standard back and forth argument that if  $K_1 = K_{r_1}$  and  $K_2 = K_{r_2}$  have property  $P_{n,s}$  for all  $n \geq 1$  and  $n \geq s \geq 0$ , then  $K_1$  and  $K_2$  are isomorphic, i.e there is a bijection  $F : \mathbb{N} \rightarrow \mathbb{N}$  such that

$$\{i, j\} \in E_{r_1} \iff \{F(i), F(j)\} \in E_{r_2}$$

**Theorem 16** *If  $K_r$  is random graph, then  $K_r$  has property  $P_{n,s}$  for all  $n \geq 1$  and  $n \geq s \geq 0$*

**Corollary 17** *If  $K_{r_1}$  and  $K_{r_2}$  are two random graphs, then  $K_{r_1}$  and  $K_{r_2}$  are isomorphic.*

## Classical Theory of Random Graphs

Consider the collection of all finite graphs on vertex set  $[n] = \{1, \dots, n\}$ .

We can put a probability measure on this set by declaring the probability of any edge occurring in such a graph is  $p$ ,  $0 < p < 1$ .

More formally for any graph  $G = ([n], E)$  we declare that the probability of  $G$  is  $p^{|E|}(1 - p)^{\binom{n}{2} - |E|}$ .

Glebskii, Kogan, Liagonkii and Talanov (1969) and, independently, Fagin (1976), proved the following theorem.

**Theorem 18** *For any fixed  $0 < p < 1$  and any first order sentence  $A$  in theory of graphs,*

$$\lim_{n \rightarrow \infty} \Pr(\{G = ([n], E) : G \models A\}) = 0 \text{ or } 1.$$

## Classical Theory of Random Graphs II

Consider the first order theory of

$$\{A : \lim_{n \rightarrow \infty} Pr(\{G = ([n], E) : G \models A\}) = 1\}.$$

It turns out that property  $P_{m,s}$  can be viewed as a first order sentence so that  $P_{m,s}$  is in this theory.

(Spencer's book, The Probabilistic Method)

**Definition** A graph  $G$  is said to have the **full level  $s$  extension property** if for any sets of vertices  $v_1, \dots, v_a$  and  $u_1, \dots, u_b$  where  $a + b \leq s$ , there is a vertex  $w$  such that for all  $i \leq a$ ,  $\{w, v_i\}$  is an edge in  $G$  and, for all  $j \leq b$ ,  $\{w, u_j\}$  is not an edge in  $G$ .

For any fixed  $0 < p < 1$ , the probability that the set of graphs  $G = ([n], E)$  such that  $G$  has the full level extension property approaches 1 as  $n$  approaches  $\infty$ .

## (Martin-Löf) Random Linear Orderings

Consider a computable dense linear ordering without end points over  $L = (\mathbb{N}, \prec)$  with universe the natural numbers.

For any real  $r = r_0 r_1 \cdots \in \{0, 1\}^{\mathbb{N}}$ , we let  $L_r = (U_r, \prec)$  where  $U_r = \{i : r_i = 1\}$ .

Then we say that  $L_r$  is **(Martin-Löf) random** if and only if  $r$  is (Martin-Löf) random.

**Theorem 19** *It  $L_r$  is random linear ordering, then  $L_r$  is dense linear ordering with out end points.*